

Superposition, Entanglement, and Quantum Computation

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Introduction - Feynman

- An N-particle quantum system can't be simulated on a classical machine whose resources don't grow exp with N.
 - *Would* be possible on a 'quantum computer'
 - *Not* a Turing machine
- Both have been proven true

Introduction, cont'd

- Quantum parallelism – quantum superposition of distinct states
 - Doesn't immediately lead to speedup
- Shor showed how info could be extracted usefully
 - Polynomial factoring algo.

Intro to



- On a classical computer, unsorted database search takes $O(n)$ time
- In 1997, Grover showed a quantum algo that takes $O(\sqrt{N})$

Superposition and entanglement

- Quantum systems can exhibit superpositions of eigensolutions
 - Not specifically quantum – classical too
- Ekert and Josza consider a multi-qubit system: apply gate U to qubits (i,j) n times
 - Quantum system: measurement in $O(n)$
 - Classical system: measurement in $O(2^n)$

Classical and quantum: a difference

- Classical waves allow superposition
 - Qubit could be represented by classical strings?
- Superposition can always be described by Cartesian product of states
- Quantum superposition may be 'entangled'
 - $\frac{1}{2}(|0\rangle + |1\rangle + |2\rangle + |3\rangle)$ can be factored
 - $\frac{1}{\sqrt{2}}(|0\rangle + |1\rangle)$ cannot be: it is entangled
- Difference is Cartesian vs. tensor products

Entanglement, cont'd

- Schroedinger says quantum entanglement is defining characteristic
- Entanglement depends on basis
 - $\frac{1}{2}(|0\rangle + |1\rangle - |2\rangle + |3\rangle)$ is entangled wrt $C_2 \times C_2$, but not wrt C_4
- State of n qubits is 2^n -dim, isomorphic to 1 particle with 2^n levels.
 - Not useful for complexity consideration, as the 1 particle requires energy resources in $O(2^n)$

Back to Grover

- Search through a phone book for name, only knowing telephone number
 - Takes $O(n)$ time classically
 - $O(\sqrt{n})$ time by Grover's algo

Basics

- There are $N = 2^L$ states labelled $S_0, S_1, S_2 \dots S_{N-1}$
 - Only one fulfills the condition C_J so that $C_J(S_J) = 1$ and $C_J(S_K) = 0, K \neq J$
- Goal is to find the solution S_J in the fewest evaluations of C_J

Grover's solution

- Start with an L-qubit register in state $|0\rangle$
- Apply an L-qubit Hadamard gate, yields an equal superposition
- Perform the following two operations on the wires, $O(\sqrt{N})$ times:

Grover's operations

- 1) Apply oracle U_J defined by:
 - $U_J |J\rangle = -|J\rangle$
 - $U_J |K\rangle = |K\rangle, K \neq J$
- 2) Apply diffusion operator D :
 - $D = H U_0 H$
 - $U_0 |0\rangle = -|0\rangle$
 - $U_0 |K\rangle = -|K\rangle, K \neq 0$

Result

- After $O(\sqrt{N})$ iterations, the outcome is the state $|J\rangle$ with high probability
- Grover explains D to be an 'inversion about the average' of the coefficients

Example

- Apply Hadamard to get
 - $|M\rangle = \frac{1}{2}(|0\rangle + |1\rangle + |2\rangle + |3\rangle)$
- Now apply the oracle U_J
 - $U_J|M\rangle = |M\rangle - 2\langle J|M\rangle |M\rangle$
- Apply Grover's diffusion operator:
 - $D U_J H|J\rangle = -|J\rangle$
 - Found in one pass!

Classical implementation

- We map each integer $0 \dots 2^L - 1$ into another integer in the same range:
 - Define L qubits to be a 'control' register $|J\rangle$ and another L to be the 'target' register $|K\rangle$
 - Let $|J\rangle \times |K\rangle \rightarrow |J\rangle \times |K \times f(J)\rangle$
 - Starting with $|K\rangle = 0$, we get $|f(J)\rangle$

Classical, cont'd

- Consider an $f(M)$ that maps an integer M to an integer $F = f(M)$ (bijective)
- Want to force init state into $|M\rangle \times |F(M)\rangle$ so that we can measure $f^{-1}(F) = M$
 - Define $W = V_f H_c$
 - Let U_f be an oracle that flips the sign of the state iff it is $|F\rangle$

An Electronic approach

- Use 2^n signal paths, one for each base state
- L -qubit Hadamard device uses op-amps with 2^n inputs and outputs
- (Description of how they used motherboards with what color LEDs here)

Hadamard implementation

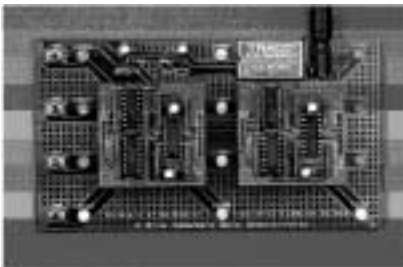
- A general L -qubit Hadamard operator can be written as a $2^L \times 2^L$ matrix
- Split each of the 2^L input signals into 2^L separate signals, each with amplitude $1/\sqrt{2^L}$
- Use an inverting op-amp for phase-shift

Electronic Hadamard



Fig. 3. Schematic diagram for the single-qubit Hadamard gate.

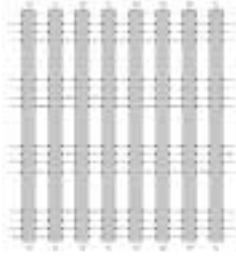
A photograph



Hadamard conclusion

- Is reversible: two applications always restores input
- Is not *physically* reversible
- Use of op-amps and resistors ensures correct operation with AC signals
- Requires 2^{2L} signals (analogous to Deutsch's 'extra universes')
- *This is just a demonstration*

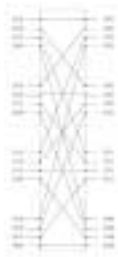
Grover schematic



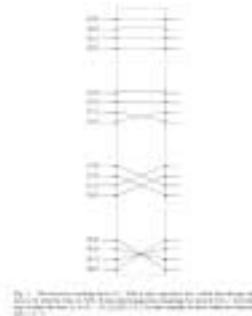
Hc schematic



T matrix



V_f for $f(l) = 3-l$



Oracle for ex. $f(l) = 2$



Conclusions

- Entanglement depends on the representation
- Their electronic implementation shows that any implementation without multi-particle entanglement requires exp. resources (refer to Ekert and Josza)

Final conclusion

- “The number of signal paths increases exponentially and makes electronic implementations of large numbers of qubits impracticable”
- Therefore, multi-particle entanglement is the key property of quantum systems that gives rise to the remarkable power of quantum computers